

Containment problem

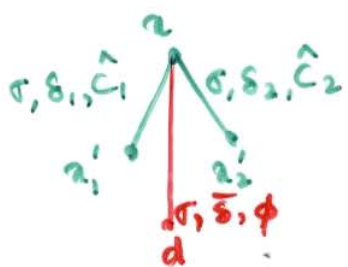
- Given G_1, G_2 , $T_m(G_1) \subseteq T_m(G_2)$?
- Containment first reduced to emptiness for deterministic G_2

$$T_m(G_1) \subseteq T_m(G_2) \iff T_m(G_1 \parallel \bar{G}_2^c) = \emptyset$$

(for nondeterministic G_2 , containment undecidable)

- determinism \iff given an activity state and an event label only one transition is enabled for a clock value

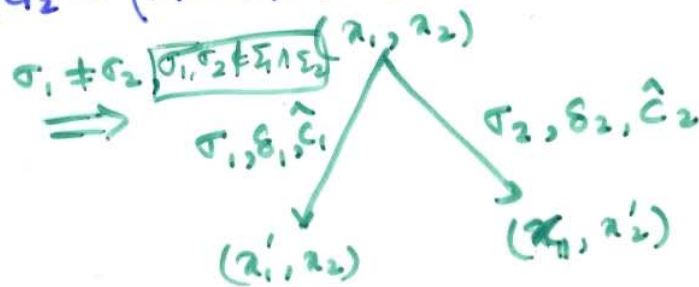
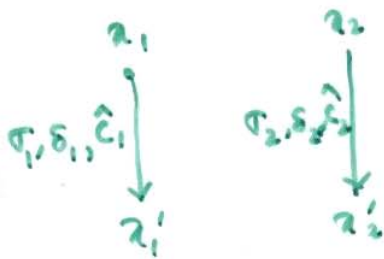
- completion: add a dump state "d" in G_2 , and
 $\forall x \in X, \sigma \in \Sigma$: add transition $(x, \sigma, \bar{\delta}, \phi, d)$, where
 $\bar{\delta} = \neg(\bigvee_i \delta_i)$, $(x, \sigma, \delta_i, \hat{c}_i, x'_i) \in E$ for each i .



determinism $\iff \delta_1 \wedge \delta_2 = \text{false}$

- Complementation: complement the markings of activity state

Composition: $G_1 \parallel G_2 = (X_1 \times X_2, \Sigma_1 \cup \Sigma_2, C_1 \cup C_2, (x_{01}, x_{02}), X_{m1} \times X_{m2}, E)$



$\sigma_1 = \sigma_2 = \sigma$
 \implies

